

Modular tensor categories in an irrational world

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Outline

1. Rationality and beyond!
2. Modularity
3. Examples
4. Why should we care?
5. Conclusions

Rationality

Conformal field theory (CFT) is quantum field theory with invariance under conformal (angle-preserving) transformations.

In two dimensions, local conformal transformations are (anti)analytic. They give rise to two commuting copies of the Virasoro algebra.

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The space of states \mathbf{H} of the CFT is thus a $Vir \oplus Vir$ -module.

More generally, \mathbf{H} is a module of two commuting copies of the symmetry algebra of the CFT, a vertex operator algebra (VOA) V .

Definition

A CFT is **rational** if \mathbf{H} is

- *semisimple* as a $V \otimes V$ -module; and
- decomposes into a *finite* number of simple $V \otimes V$ -modules.

Examples and non-examples

1. The **Ising model** is rational with V being the simple Virasoro VOA of central charge $\frac{1}{2}$ and

$$H = (L_0 \otimes L_0) \oplus (L_{1/16} \otimes L_{1/16}) \oplus (L_{1/2} \otimes L_{1/2}),$$

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2. The **free boson** is not rational with V being the Heisenberg VOA of central charge 1 and

$$H = \int_{\mathbb{R}}^{\oplus} (F_p \otimes F_p) dp,$$

where F_p is the Fock space of charge p . Whilst H is semisimple, it is composed of an uncountably infinite number of simples.

3. The **triplet model** is not rational. It has four simple V -modules — W_0 , W_1 , $W_{-1/8}$ and $W_{3/8}$ — and the state space decomposes as

$$H = (W_{-1/8} \otimes W_{-1/8}) \oplus (W_{3/8} \otimes W_{3/8}) \oplus \left[\begin{array}{l} \text{an indecomposable} \\ \text{with 8 composition} \\ \text{factors built from} \\ W_0 \text{ and } W_1 \end{array} \right].$$

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4. The **singlet model** is not rational either. It has an uncountable infinity of simple V -modules E_λ , $\lambda \in \mathbb{C} \setminus \mathbb{Z}$, and M_r , $r \in \mathbb{Z}$, with

$$H = \int_{\mathbb{C} \setminus \mathbb{Z}}^{\oplus} (E_\lambda \otimes E_\lambda) d\lambda \oplus \left[\begin{array}{l} \text{an indecomposable with} \\ \text{a countable infinity of} \\ \text{composition factors built} \\ \text{from the } M_r \end{array} \right].$$

It is neither finite nor semisimple [Wang'97].

Message of the day

Highly opinionated message.

From the modular perspective, the difficulty of these four cases is

$$1. < 2. < 4. < 3.$$

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Starting with 3. is silly...

Modularity

Recall that the (unqualified) modular group is

$$\begin{aligned} \mathrm{SL}_2(\mathbb{Z}) &= \left\{ \begin{pmatrix} a & b \\ c & d \end{pmatrix} : a, b, c, d \in \mathbb{Z}, ad - bc = 1 \right\} \\ &= \langle S, T : S^2 = (ST)^3 = S^{-2} \rangle. \end{aligned}$$

It is common to identify S and T with the matrices

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A **consistency condition** of CFT is that the partition function (character of \mathbf{H}) is invariant under the natural action of the modular group.

Verlinde noticed/conjectured that the fusion coefficients \mathcal{N}_{ij}^k , given by

$$M_i \times M_j \simeq \bigoplus_k \mathcal{N}_{ij}^k M_k,$$

are determined by the modular S-transforms of the characters,

$$\text{ch}[M_i] \xrightarrow{S} \sum_j S_{ij} \text{ch}[M_j],$$

via the celebrated (and highly non-trivial) **Verlinde formula**:

$$\mathcal{N}_{ij}^k = \sum_{\ell} \frac{S_{i\ell} S_{j\ell} S_{k\ell}^*}{S_{1\ell}} \quad (M_1 \text{ is the VOA}).$$

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[Moore–Seiberg'89] explained that this (essentially) follows from other consistency conditions of CFTs.

This led [Turaev'94] to define the notion of a **modular fusion category** (MFC), which means (among other things) that:

- the module category is finite + semisimple (+ ...);
- the module category admits a tensor product (**fusion**);
- there is a “nice” action of the modular group $SL_2(\mathbb{Z})$; and
- the fusion product and the modular action are related by Verlinde.

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Huang proved that rational CFTs are modular. But, which generalisation of an MFC captures the modularity of non-rational CFTs?

Again with the examples...

1. The Ising model's three simples have characters given by

$$\begin{aligned}\text{ch}[\mathbf{L}_0] &= \frac{1}{2} \left(\sqrt{\frac{\vartheta_3(0,\tau)}{\eta(\tau)}} + \sqrt{\frac{\vartheta_4(0,\tau)}{\eta(\tau)}} \right), \\ \text{ch}[\mathbf{L}_{1/2}] &= \frac{1}{2} \left(\sqrt{\frac{\vartheta_3(0,\tau)}{\eta(\tau)}} - \sqrt{\frac{\vartheta_4(0,\tau)}{\eta(\tau)}} \right), \\ \text{ch}[\mathbf{L}_{1/16}] &= \sqrt{\frac{\vartheta_2(0,\tau)}{2\eta(\tau)}}.\end{aligned}$$

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This CFT is modular with S- and T-matrices given by

$$\mathcal{S} = \frac{1}{2} \begin{pmatrix} 1 & \sqrt{2} & 1 \\ \sqrt{2} & 0 & -\sqrt{2} \\ 1 & -\sqrt{2} & 1 \end{pmatrix}, \quad \mathcal{T} = e^{-i\pi/24} \begin{pmatrix} 1 & 0 & 0 \\ 0 & e^{i\pi/8} & 0 \\ 0 & 0 & -1 \end{pmatrix}.$$

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The Verlinde formula gives non-negative integers that agree with the fusion coefficients, as required.

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$$\mathcal{S}(p, p') = e^{-2\pi i p p'}, \quad \mathcal{T}(p, p') = e^{-\pi i/12} e^{\pi i p^2} \delta(p - p').$$

Note that \mathcal{S} is just a Fourier transform!

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Even though the free boson is not rational, the Verlinde-like formula (with $\sum \rightarrow \int$) still gives the fusion coefficients.

3. The triplet model has four simples whose characters are

$$\begin{aligned} \text{ch}[W_0] &= \frac{1}{2} \left(\frac{\vartheta_2(0, \tau)}{2\eta(\tau)} + \eta(\tau)^2 \right), & \text{ch}[W_{-1/8}] &= \frac{\vartheta_3(0, \tau) + \vartheta_4(0, \tau)}{2\eta(\tau)}, \\ \text{ch}[W_1] &= \frac{1}{2} \left(\frac{\vartheta_2(0, \tau)}{2\eta(\tau)} - \eta(\tau)^2 \right), & \text{ch}[W_{3/8}] &= \frac{\vartheta_3(0, \tau) - \vartheta_4(0, \tau)}{2\eta(\tau)}. \end{aligned}$$

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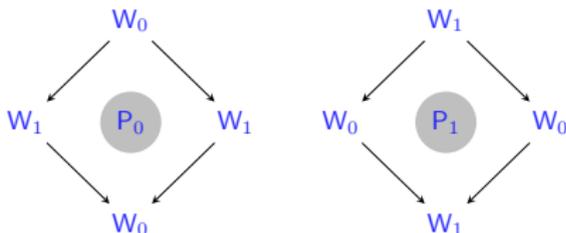
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W_0 and W_1 harbour the non-semisimplicity of the CFT. In particular, their projective covers are reducible but indecomposable.



We do have an action of $SL_2(\mathbb{Z})$ on the **projective** characters, but they are not linearly independent: $\text{ch}[P_0] = \text{ch}[P_1]$.

This gives a modular-invariant partition function:

$$\frac{1}{2} \left| \text{ch}[P] \right|^2 + \left| \text{ch}[W_{-1/8}] \right|^2 + \left| \text{ch}[W_{3/8}] \right|^2.$$

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However, the S-matrix is no good for Verlinde interpretations:

$$\mathcal{S} = \begin{pmatrix} 0 & 1 & -1 \\ \frac{1}{2} & \frac{1}{2} & \frac{1}{2} \\ -\frac{1}{2} & \frac{1}{2} & \frac{1}{2} \end{pmatrix}.$$

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On the other hand, the projectives carry a non-semisimple action of L_0 , so one can introduce a pseudotrace to see the non-trivial Jordan blocks.

Here, this augments the characters by the **pseudocharacter** $-i\tau \eta(\tau)^2$. Together, they span a 5-dimensional $SL_2(\mathbb{Z})$ -module (but it is still no good for Verlinde games).

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Moreover, the characters of the **atypical modules** are modular “almost everywhere” if we allow certain infinite-linear combinations:

$$\text{ch}[\mathbf{M}_r] = \sum_{j \geq 0} (-1)^j \text{ch}[\mathbf{E}_{r+j+1}]$$

$$\Rightarrow \mathcal{S}(r, \lambda') = \frac{e^{-2\pi i r(\lambda' - \frac{1}{2})}}{2 \cos[\pi(\lambda' - \frac{1}{2})]}, \quad (r \in \mathbb{Z}, \lambda' \in \mathbb{R}).$$

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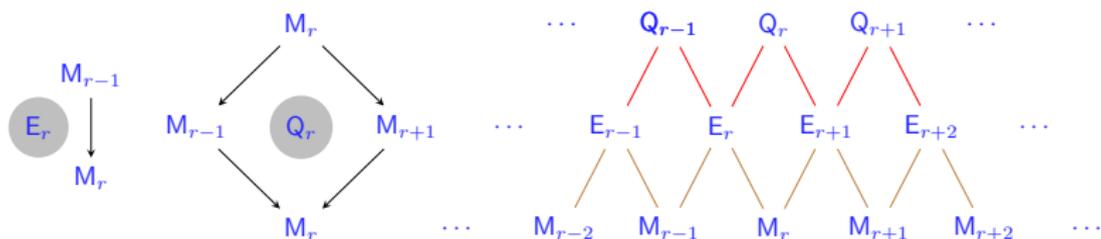
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Most importantly (and surprisingly), the Verlinde formula works!

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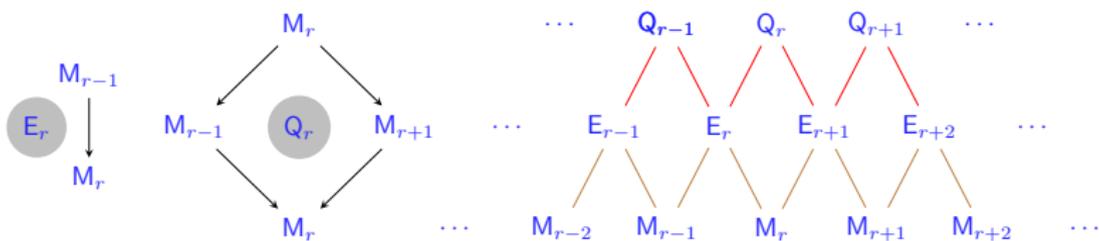
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Q_r is the projective cover of M_r — it is a **tilting** module.

Moreover, these modules exhibit **BGG duality**:

$$[Q_r : E_s] = [E_s : M_r].$$

In fancy terms, the singlet category is a **BGG highest-weight category**.

Why should we care?

These examples make manifest the claim that finite + non-semisimple is harder than non-finite + non-semisimple. We thus expect good modularity when the CFT is “semisimple almost everywhere”.

Another controversial opinion.

Imposing finiteness on modularity is as silly as studying theta functions without Poisson resummation and Fourier analysis.

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This observation is the basis for the **standard module formalism** of [Creutzig–DR’13, DR–Wood’14].

- The trivial cases are the semisimple ones (1. & 2.): simple = standard = projective.
- In the non-finite + non-semisimple case (4.), all examples studied show a BGG/highest-weight structure like the singlet.
- It does not apply (directly) to the finite + non-semisimple case (3.), because the non-semisimplicity occupies a set of positive measure.

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Unfortunately, we’re rubbish at constructing examples of type 3.

A better quest is thus to try to categorically characterise the standard modularity of non-finite + non-semisimple examples (*cf.* MFCs).

These are **modular BGG categories**.¹

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Part of this requires developing tools to compute correlation functions for these examples (and so extract categorical data).

One can then use this to inform the modular behaviour of finite + non-semisimple examples like the triplet.

We also need to construct more examples of this type, so as to determine which properties are natural/general.

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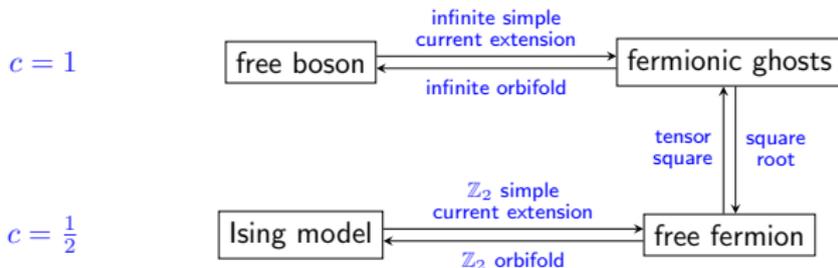
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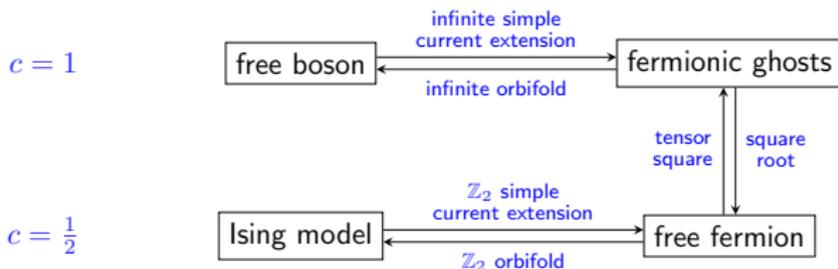


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In particular, infinite simple-current extensions often do this.

Example:



Many rational CFTs (eg. some W-algebras) are only known because they have been constructed from non-rational CFTs.

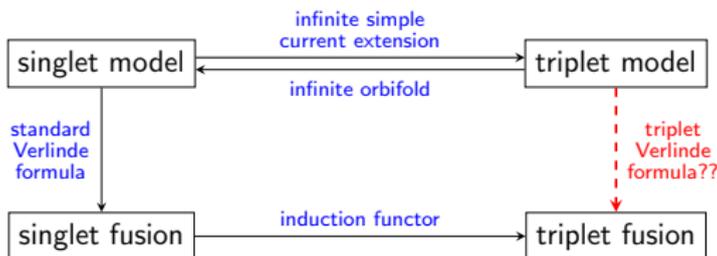
Fixing the triplet

There have been several articles proposing Verlinde-like formulae for the triplet model, eg. [Fuchs-Hwang-Semikhatov-Tipunin, Gaberdiel-Runkel, Gainutdinov-Runkel, Creutzig-Gannon].

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However, the triplet is related to the singlet by an infinite-order simple-current extension [Creutzig-DR'13]. The easiest approach to analysing triplet modularity exploits this fact [Melville-DR'15].



A Verlinde formula for the triplet

As above, compute the singlet fusion rules using the standard Verlinde formula for the singlet, use induction to recover triplet fusion coefficients and then try to rewrite everything in terms of triplet modular data.

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This is delicate (but possible). The result is a modified Verlinde formula that takes into account the partition of simples into:

$$\text{atypical: } \{W_0, W_1\} \quad \text{typical: } \{W_{-1/8}, W_{3/8}\}.$$

The triplet Verlinde formula is:

$$\mathcal{N}_{ij}^k = \sum_{\ell \in \text{atyp.}} \frac{S_{i\ell} S_{j\ell} S_{\ell k}^{-1}}{S_{1\ell}} + \delta_{ijk} \sum_{\ell \in \text{atyp.}} \frac{S_{i\ell} S_{j\ell} S_{\ell k}^{-1}}{S_{1\ell}},$$

where $\delta_{ijk} = 1$ if only $i \in \text{atyp.}$, only $j \in \text{atyp.}$ or $i, j, k \in \text{atyp.}$, and $\delta_{ijk} = 0$ otherwise.

Note that the S-matrix is that of the triplet **with the factors of τ !**

Conclusions

- Modularity is subtle for non-semisimple CFTs, eg. the triplet. We do not have an $SL_2(\mathbb{Z})$ -action on the characters in general and so the Verlinde formula does not seem to make sense.

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Final inflammatory opinion.

We should stop proposing finite generalisations of modular fusion categories and instead study modular BGG categories.